# Optimizing Data Plane Programs for the Network

<u>Johannes Krude</u><sup>1</sup>, Matthias Eichholz<sup>2</sup>, Maximilian Winck<sup>1</sup>, Klaus Wehrle<sup>1</sup>, Mira Mezini<sup>2</sup>



<sup>1</sup>RWTH Aachen University, <sup>2</sup>Technische Universität Darmstadt

https://comsys.rwth-aachen.de/

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#### **Data Plane Programs**

- Shift from fixed to programmable data planes
  - P4 on e.g., Barefoot Tofino
  - ▶ eBPF on NPUs, e.g., Netronome NFP-6000
  - eBPF with XDP e.g., Intel Core-i7
- Programs are executed for each packet
- Limited per-packet time budget
  - e.g., 67.2 ns at 10GbE
- Limited program size
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#### Our proposal

Integrating the network topology into program optimization







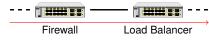


#### **Firewall**

- Drops non IPv6 TCP
  - Inserts fixed sized IPv6 options
    - e.g., Authentication Header

- 1. Parse Ethernet header
- 2. Parse IPv6 header
- 3. Loop: Parse IPv6 option
- 4. Extract IPv6 dst & TCP dst
- 5. Lookup & redirect to backend





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Parts of a data plane program are redundant if equivalent functionality was already applied to packets before reaching this program.

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# This example, XDP BPF Core i7-870

- $\Rightarrow$  **2.3**× packet rate (2.5 $\rightarrow$ 5.8 M pkt/s)
- $\Rightarrow$  **12%** program size(7.6 $\rightarrow$ 0.9 KiB)

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if (firewall(&packet) != DROP)
load_balancer(&packet);
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Krude et al.

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Problem: Conventional tools (e.g., Ilvm) only output the combined

optimized program

**Solution:** Instead use exhaustive symbolic execution

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- Replace dead branch by jump
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- Run conventional compiler passes \_\_\_
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# 3. Optimization

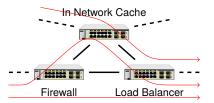
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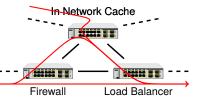
# Runtime to optimize the cilium.io Load Balancer in our example

• 30 s for all steps including 25 s of symbolic execution



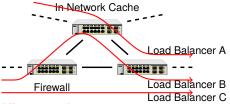


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  - A single optimized variant valid for all paths
  - Different optimized variants for different paths

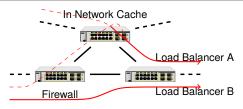




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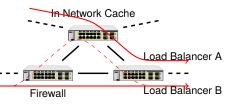


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  - Only consider neighbours
  - Leave out irrelevant intermediate nodes



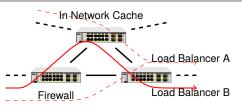


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- Selectively applying optimizations
  - Optimize slow hotspots
  - Identify program chains with high optimization potential
  - Run some optimization whenever controller is idle





#### Summary

- Our proposal: Integrating the network topology into program optimization
- High optimization potential on example with real-world program
  - 2.3× packet rate increase
  - 88% program size reduction
- Still much to do
  - Improving optimization time
  - Considering multiple network paths
  - Evaluating on realistic networks

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# **Questions?**

